

# The Unconstrained Binary Quadratic Programming Problem: A Survey

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## Abstract:

In recent years the unconstrained binary quadratic program (UBQP) has grown in importance in the field of combinatorial optimization due to its application potential and its computational challenge. Research on UBQP has generated a wide range of solution techniques for this basic model that encompasses a rich collection of problem types. In this paper we survey the literature on this important model, providing an overview of the applications and solution methods.

**Keywords:** Unconstrained binary quadratic programs, combinatorial optimization, metaheuristics

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7 **1.0 Introduction:**

8 The Unconstrained Binary Quadratic Programming (UBQP) problem is defined by

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$$\min x^t Qx$$
  
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$$s.t. x \in S$$
  
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13 where  $S$  represents the binary discrete set  $\{0,1\}^n$  or  $\{-1,1\}^n$  and  $Q$  is an  $n$ -by- $n$  square, symmetric  
14 matrix of coefficients. This simple model is notable for embracing a remarkable range of  
15 applications in combinatorial optimization. For example, the use of this model for representing  
16 and solving optimization problems on graphs, facility locations problems, resources allocation  
17 problems, clustering problems, set partitioning problems, various forms of assignment problems,  
18 sequencing and ordering problems, and many others have been reported in the literature.  
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25 Even more remarkable is the fact that, once given a UBQP formulation, these problems  
26 can be solved by a UBQP method which is not specialized to exploit the problem domain of any  
27 individual class of problems, to yield solutions whose quality in many cases rivals or even  
28 surpasses the quality of the solutions produced by the best specialized methods, while achieving  
29 this outcome with an efficiency that likewise rivals or surpasses the efficiency of leading  
30 specialized methods.  
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37 In this paper we survey the literature on UBQP, both its applications and solution  
38 methods. While many important constrained nonlinear binary models have been reported in the  
39 literature over the years, we focus our attention here on models that naturally occur in the form  
40 of an unconstrained quadratic binary model and those that have been re-cast into the form of  
41 UBQP. The paper is organized as follows. In section 2 we survey the range of applications that  
42 have been reported in the literature. Section 3 then presents a survey of the solution  
43 methodologies reported in the literature for solving UBQP. Section 4 highlights key theoretical  
44 work and this is followed by section 5 which wraps up the paper with our summary and  
45 conclusions.  
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4 **2.0 Applications:**  
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6 Some reported applications appear naturally in the form of UBQP while others are “re-  
7 cast” into the UBQP form by employing various transformations. In the sub-sections below we  
8 examine these different categories of applications in turn. Within sub-sections, we present  
9 applications in the chronological order in which they appeared in the literature to give the reader  
10 a sense of when certain topics were addressed appeared in print, as well as progress made and  
11 trends in solution methodologies.  
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19 **2.1 Natural UBQP problems/applications**  
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21 The literature on UBQP goes back to the 1960s where the topics of pseudo-boolean  
22 functions and binary quadratic optimization were introduced by Hammer and Rudeanu (1968).  
23 Early papers *related* to UBQP concern applications in finance (Laughunn (1970)), project  
24 selection (Rhys(1970)), cluster analysis (Rao (1971)), economic analysis (Hammer and Shliffer  
25 (1971)), traffic management (Witzgall(1975)) and computer aided design (Krarup and Pruzan  
26 (1978)). While these applications actually take the form of constrained quadratic binary  
27 programs, they are mentioned here due to their historical role in fostering an interest in quadratic  
28 binary applications and also because several allow special cases that are precisely in the form of  
29 UBQP.  
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37 More recently many interesting applications that are expressed naturally in the form of  
38 UBQP have appeared in various papers. Barahona, Grotschell, Junger and Reinelt (1988)  
39 formulate and solve the problem of finding ground states of spin glasses with exterior magnetic  
40 fields, an important problem in physics, as an instance of UBQP. Computational results reveal  
41 that the model produces high quality solutions to spin glass problems of realistic size in  
42 reasonable amounts of computation time using 1980s technology.  
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48 Hansen and Jaumard (1990), in their work on the satisfiability problem, report their  
49 experience using the UBQP model as an approach for representing and solving small to medium  
50 sized Max 2-sat problems. Computational studies validated the attractiveness of this approach to  
51 the Max 2-sat problem in terms of quickly producing high quality solutions.  
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55 Boros and Hammer (1991) discuss the use of UBQP as an approach for modeling the  
56 Max-Cut problem. Their paper highlights the relationship between UBQP, Max-cut, Max 2-sat,  
57 and the Weighted Signed Graph Problem. The authors also present a discussion of valid  
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4 inequalities and facets for polyhedra that provide the basis for further computational and  
5 theoretical work.  
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8 Alidaee, Kochenberger and Ahmadian (1994) discuss two machine scheduling problems  
9 in the context of UBQP: (1) scheduling  $n$  jobs on a single machine to minimize total weighted  
10 earliness and tardiness, and (2) scheduling  $n$  jobs on two parallel identical processors to  
11 minimize weighted mean flow time. In each case, the authors show how the problems can be  
12 modeled in a straight-forward manner as an instance of UBQP.  
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17 Pardalos and Xue (1994) indicate how the maximum clique problem can be modeled as  
18 an instance of UBQP. The authors also discuss the relationship between the maximum clique  
19 problem, the maximum independent set problem, and the vertex cover problem, indicating how  
20 each can be represented by UBQP. Finally, the authors provide a survey of solution methods for  
21 the maximum clique problem.  
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27 De Simone, Diehl, Junger, Mutzel, Reinelt, and Rinaldi (1995), as in the earlier 1988  
28 paper by Barahona, Grotschell, Junger and Reinelt, adopt the UBQP model as a representation  
29 for the problem of finding ground states for the spin glass problem. In this 1995 paper the  
30 authors use the UBQP model to compute exact ground states for Ising spin glasses on 2-  
31 dimensional grids with periodic boundary interactions, Gaussian bond distributions, and an  
32 exterior magnetic field. Preliminary experiments with a branch and cut algorithm for optimizing  
33 the UBQP form of the problem proved very promising, quickly producing high quality solutions  
34 to large spin glass instances.  
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41 Bomze, Budinich, Pardalos, and Pelillo (1999) discuss the maximum clique (MC)  
42 problem and how it can be modeled in a variety of ways including a representation in  
43 terms of UBQP. The authors provide a very broad and in depth discussion of a variety of  
44 applications and of both exact and heuristic solution methods for the MC problem.  
45 Computational experience with various solution approaches to the MC problem is also presented.  
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51 Iasemidis, Pardalos, Sackellares, and Shiau (2001) discuss the use of the UBQP model as  
52 part of a process employed to predict the arrival of epileptic seizures. The entrainment between  
53 two brain sites can be quantified from measures of electrical activity (EEG) of the brain. The  
54 UBQP model was successful in identifying the most entrained sites leading to the optimal  
55 location of electrode sites. In clinical trials this procedure was successful in predicting epileptic  
56 seizures 20-40 minutes in advance of their occurrence.  
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4 Alidaee, Glover, Kochenberger, and Rego (2005) discuss the number partitioning  
5 problem where the objective is to assign numbers to subsets such that the sums of the numbers in  
6 each subset are as close as possible to one another. The authors show that the  $n = 2$  subset case  
7 can be modeled as an instance of UBQP and that problems with  $n > 2$  can be modeled as a  
8 constrained version of UBQP. Extensions of the basic model along with computational  
9 experience for the  $n=2$  case are presented indicating the attractiveness of the approach.  
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15 Kochenberger, Glover, Alidaee, and Lewis (2005) discuss their experience with adopting  
16 the UBQP model to represent and solve max 2-sat problems. Expanding the computational  
17 scope reported earlier by Hansen and Jaumard (1990) on UBQP and the Max 2-sat problem, they  
18 offer extensive computational experience on very large test problems with up to 1000 variables  
19 and more than 10,000 clauses. Employing a basic form of tabu search to solve the UBQP  
20 instances, best known solutions to most test problems were found in a few seconds of  
21 computation time.  
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28 Neven, Rose and Macready (2008) discuss the use of quantum adiabatic algorithms,  
29 which represent new approaches to NP-hard combinatorial problems, for solving the image  
30 recognition problem. The authors indicate how the pattern recognition problem of deciding  
31 whether two images contain the same object can be modeled as an instance of UBQP, which they  
32 show is the general input format required by D-Wave superconducting quantum AQC  
33 processors. Computational experience was not reported.  
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39 Pajouh, Balasundaram, and Prokepyev (2013) discuss the use of the UBQP model for  
40 representing the maximal independent set problem. The authors present an analysis of local  
41 maxima properties along with relations between continuous local maxima of the quadratic  
42 formulation and the binary local maxima in the Hamming distance 1 and 2 neighborhoods.  
43 These results are then used to construct effective local search algorithms for the maximum  
44 independent set problem.  
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50 Kochenberger, Hao, Lu, Glover and Wang (2013) discuss the Max Cut problem and how  
51 the UBQP model can be effectively used to model and solve large scale instances. Using a tabu  
52 search algorithm, extensive computational testing is reported on problems with up to 10,000  
53 variables. Comparisons with other solution methods from the literature for the max cut problem  
54 are provided, indicating the attractiveness of the UBQP/Tabu Search approach.  
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## 2.2 UBQP via reformulation

The applications of the previous section illustrate the widespread usefulness of the UBQP model. The actual applicability of UBQP, however, is greatly extended due to re-formulation procedures that re-cast a *constrained* problem into an equivalent unconstrained binary quadratic model. Many re-formulations are accomplished by including quadratic infeasibility penalties in the objective function as an alternative to explicitly imposing constraints. In this manner a constrained model can be re-cast into the form of UBQP. In fact, any linear or quadratic problem with linear constraints and bounded integer variables can in principle be re-formulated as UBQP using quadratic penalties.

For several simple constraints, appropriate quadratic penalties are known in advance and can be used straight away. Examples of such penalties are given in Table 1 below where  $P$  is a large positive scalar.

Classical Constraint	Equivalent Penalty
$x + y \leq 1$	$P(xy)$
$x + y \geq 1$	$P(1 - x - y + xy)$
$x + y = 1$	$P(1 - x - y + 2xy)$
$x \leq y$	$P(x - xy)$
$x_1 + x_2 + x_3 \leq 1$	$P(x_1x_2 + x_1x_3 + x_2x_3)$

Table 1: Illustrative Known Penalties

Note that the penalty term in each case is zero if the associated constraint is satisfied, and otherwise the penalty is positive. These penalties, then, can be directly employed as an alternative to explicitly introducing the original constraints. For general constraints, however, appropriate penalty functions are not known in advance and need to be “discovered.” A simple procedure (see for instance Hammer & Rudeanu (1968); Hansen (1979); Hansen, Jaumard, & Mathon (1993); and Boros & Hammer( 2002) for finding an appropriate penalty for any linear constraint is given as follows:

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4 Consider the general constrained problem of the form

$$5 \quad \min x_0 = xQx$$

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$$9 \quad \text{s.t.} \quad Ax = b, \quad x \text{ binary} \quad (1)$$

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15 This model accommodates both quadratic and linear objective functions since the linear case  
16 results when Q is a diagonal matrix (observing that  $x_j^2 = x_j$  when  $x_j$  is a 0-1 variable). Under the  
17 assumption that A and b have integer components, problems with inequality constraints can also  
18 be put in this form by representing their bounded slack variables by a binary expansion. These  
19 constrained quadratic optimization models are converted into equivalent UQP models by adding  
20 a quadratic infeasibility penalty function to the objective function in place of explicitly imposing  
21 the constraints  $Ax = b$ .  
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28 Specifically, for a positive scalar P:

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$$32 \quad x_0 = xQx + P(Ax - b)^t (Ax - b)$$

$$33 \quad = xQx + xDx + c$$

$$34 \quad = x\hat{Q}x + c$$

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$$37$$

38 where the matrix D and the additive constant c result directly from the matrix multiplication  
39 indicated. Dropping the additive constant, the equivalent unconstrained version of the  
40 constrained problem becomes  
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$$44 \quad UBQP : \min x\hat{Q}x, \quad x \text{ binary} \quad (2)$$

$$45$$

$$46$$

47 A suitable choice of the penalty scalar P can always be chosen so that the optimal solution to  
48 UBQP is the optimal solution to the original constrained problem. For ease of reference, the  
49 preceding procedure that transforms (1) into (2) will be called *Transformation # 1*.  
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53 Transformation #1 can be used in cases where an appropriate quadratic penalty function  
54 isn't known in advance. In certain special cases, as mentioned earlier, appropriate penalties are  
55 known and can be directly employed. One particularly important case that arises in many  
56 constrained combinatorial problems is:  
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$$x_j + x_k \leq 1 \quad (3)$$

denoting a situation where a pair of binary choices are available and we must preclude choosing both. As shown in the preceding table, an equivalent quadratic penalty for this situation is simply

$$Px_jx_k \quad (4)$$

Due to the frequency with which the constraint of (3) appears in many important applications we single it out for special attention and refer to the penalty of (4) as an alternative to the constraint of (3) as *Transformation # 2*. Many of the applications that follow were originally modeled as constrained 0/1 models and were recast into the form of UBQP by using Transformation 1 and/or 2.

Other paths to reformulation exist as well. Often a well-chosen change of variable can result in transforming a constrained model into the form of UBQP. This is particularly important in the context of certain optimization problems on graphs where binary variables denoting whether or not an edge is chosen can be replaced by the product of the two associated binary node variables. In making such a substitution, we go from an “edge-oriented” model to a “node-oriented” model. This typically results in a much smaller model in terms of both the number of variables and the number of constraints.

The clique partitioning problem affords a good example for illustrating this approach. The standard integer programming (IP) formulation (see for example Oosten, Rutten & Spijksma (2001)) for clique partitioning is:

$$\max \sum_{(i,j) \in E} w_{ij}x_{ij}$$

s.t.

$$x_{ij} + x_{ir} - x_{jr} \leq 1 \quad \forall \text{ all distinct } i, j, r \in V$$

$$x_{ij} \in \{0,1\} \text{ for all } \{i, j\} \in E$$

The variable  $x_{ij}$  is equal to 1 if the edge (i,j) is in the partition and is equal to 0 otherwise. The coefficient  $w_{ij}$  is the weight of the edge (i,j) in the graph.



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4 An alternative model results by changing from edge-based variables to node-based  
5 variables. For this new model we add artificial edges as needed to produce a complete graph and  
6 denote an upper bound on the number of cliques to be formed by  $K_{\max}$ . Then, letting  $x_{ik} = 1$  if  
7 node  $i$  is assigned to clique  $k$  and  $x_{ik} = 0$  otherwise, an equivalent model is:  
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$$10 \quad \max \sum_{i=1}^{n-1} \sum_{j=i+1}^n w_{ij} \sum_{k=1}^{K_{\max}} x_{ik} x_{jk}$$

$$11 \quad s.t.$$

$$12 \quad \sum_{k=1}^{K_{\max}} x_{ik} = 1 \quad \text{for } i = 1, n$$

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18 In this formulation  $n$  is the number of nodes in the graph and  $w_{ij}$  again denotes the weight of  
19 edge  $(i,j)$ . This model is much smaller than the standard IP in terms of both number of variables  
20 and the number of constraints. Note also that it is of the form:  
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$$23 \quad \max x'Qx$$

$$24 \quad s.t.$$

$$25 \quad Ax = b \quad x \text{ binary}$$

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35 which can be re-cast into the form of UBQP using Transformation #1.  
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### 37 38 2.3 Specific Application Instances

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40 Each of the applications presented below were originally modeled as a constrained  
41 combinatorial problem and then re-cast into the form of UBQP. Once in this unified form the  
42 problems were successfully solved by various heuristic means.  
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48 Lewis, Alidaee, and Kochenberger (2005) address the problem of assigning tasks to  
49 processors in a distributed, multitasking computer architecture such that the sum of the resultant  
50 task completion costs and inter-task communication costs are minimized. The standard model  
51 for this problem is a constrained quadratic optimization model in binary variables with  
52 constraints ensuring that each task gets assigned to one of the processors available. The authors  
53 employ Transformation # 1 to re-cast this model into the form of UBQP which in turn is solved  
54 with a basic tabu search heuristic. Computational experience with large-scale instances  
55 highlights the attractiveness of this approach.  
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4 Kochenberger, Glover and Alidaee (2005) discuss the classic vertex coloring problem  
5 and how it can be effectively modeled and solved in the form of UBQP. A standard  
6 representation of the K-Coloring problem consists of two categories of constraints, one ensuring  
7 that each node gets a color, and the other ensuring that adjacent nodes receive different colors.  
8 The authors use Transformation # 1 on the first set of constraints and Transformation # 2 on the  
9 second to produce a UBQP representation of the problem. Computational experience applying a  
10 tabu search method to standard test problems from the literature indicates that this approach is  
11 very competitive with, and often superior to, specialized methods for vertex coloring.  
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20 Similar reformulations, using transformations #1 and/or #2 have been reported for other  
21 well-known combinatorial problems. Kochenberger, Glover, Alidaee, and Wang (2005) examine  
22 the use of UBQP as a tool for clustering microarray data into groups with high degrees of  
23 similarity. Wang, Alidaee, Glover, and Kochenberger (2006) discuss the problem of grouping  
24 machines and parts together in a flexible manufacturing system in a manner that facilitates  
25 economies in time and cost. Kochenberger, Alidaee, Glover, and Wang (2007) discuss the use of  
26 UBQP as a tool for modeling and solving the generalized independent set (GIS) problem. In  
27 each case the original model was re-cast into the form of UBQP and successfully solved in this  
28 new form.  
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37 In addition, Lewis, Kochenberger, and Alidaee (2008) discuss the classic set partitioning  
38 (SP) problem and how the UBQP framework can be utilized for modeling and solving this  
39 important class of problems. Computational experience using a basic tabu search heuristic on  
40 problems with up to 15,000 variables and 5,000 rows and various densities is presented with  
41 comparisons drawn with CPLEX. Also in 2008, Alidaee, Kochenberger, Lewis, Lewis, and  
42 Wang (2008) discuss the use of the UBQP model for representing and solving the well-known  
43 set packing problem. Favorable computational experience with a wide variety of set packing  
44 problems with up to 2,000 variables and 10,000 constraints is reported.  
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52 Lewis, Alidaee, Glover, and Kochenberger (2009) discuss the Linear Ordering (LO)  
53 problem and how it can be modeled and solved as an instance of UBQP. The standard model in  
54 the literature for LO is a large 0-1 linear program with many constraints. For instance, a model  
55 designed to order 150 items would have, in the classic model, more than one million constraints.  
56 Rather than use the general procedure of Transformation # 1, the authors show how to easily re-  
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4 cast the constrained model into the unconstrained form of UBQP by using a special quadratic  
5 penalty that is uniquely suitable for the problem at hand. Computational experience with both  
6 medium and large sized test problems reveals the effectiveness of this approach.  
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10 Douiri and Elberoussi (2012) discuss the Sum Coloring Problem which generalizes the  
11 classical vertex coloring problem by seeking a valid coloring of vertices such that the sum of the  
12 colors assigned to all vertices is minimized. The transformation to UBQP is accomplished by  
13 using Transformation # 1 on the constraints that ensure each vertex gets a color and  
14 Transformation # 2 on the constraints that require adjacent nodes to have different colors.  
15 Computational experience with the resulting UBQP model was carried out using a genetic  
16 algorithm. Results obtained on a variety of standard test problems illustrate the attractiveness of  
17 this approach.  
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26 Wang and Xu (2013) discuss another variant of the classical vertex coloring problem,  
27 called the Robust Graph Coloring Problem (RGCP), where for a given feasible coloring, a  
28 penalty is incurred for each non-adjacent vertices that have the same color assigned. The  
29 optimization problem is to determine a feasible coloring that minimizes the sum of the penalties  
30 associated with the edges in the complementary edge set with endpoints that are assigned the  
31 same color. As with the previously discussed coloring problems, the transformation to UBQP is  
32 carried out using a combination of Transformation #1 and Transformation #2. Computational  
33 experience with several variations of a genetic algorithm on a set of test problems illustrates the  
34 effectiveness of the UBQP approach for modeling and solving RGCP.  
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44 Lewis, Kochenberger, Wang, and Glover (2013) discuss the use of UBQP for modeling  
45 and solving the Generalized Vertex Covering Problem (GVCP). GVCP generalizes the  
46 minimum weight vertex covering problem by employing a three tier cost structure for each edge  
47 and charging a cost depending on whether one, both or neither end point of a given edge is  
48 covered by the subset chosen. The optimization problem is to find the subset of nodes that  
49 minimizes the sum of both the node and edge costs. The model presented previously in the  
50 literature for GVCP is a large 0-1 linear program with a binary variable for each node in the  
51 graph and two binary variables for each edge. The authors here show how GVCP can readily be  
52 formulated as UBQP by employing a simple change of variable such that all edge variables and  
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4 all constraints are eliminated. Computational experience comparing the original linear model  
5 and the equivalent UBQP model illustrates the superiority of UBQP for this class of problems.  
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### 8 9 **3.0 Solution Methods**

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11 While a few special cases of UBQP are polynomially solvable (see for instance Picard  
12 (1976), Barahona (1986), Pardalos and Jha (1991)), UBQP in general is an NP-hard problem  
13 (see Pardalos and Jha (1992)) and for all but small to moderate sized problems, heuristic methods  
14 are required to produce good solutions in a reasonable amount of computer time. Nonetheless,  
15 there is a sizable literature on exact methods for UBQP. In the sub-sections below we first  
16 survey the exact methods that have appeared in the literature followed by the heuristic methods  
17 described in the literature for solving UBQP.  
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#### 25 **3.1 Exact Methods**

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27 The literature on exact methods for UBQP introduces a variety of algorithms, each with  
28 the virtue of terminating, given enough time and memory, with a globally optimal solution.  
29 Most approaches involve a tree search of a general branch-and-bound nature but other methods  
30 exist as well. In this section we survey, in chronological order, the prominent methods reported  
31 over the past thirty-plus years.  
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38 Gulati, Gupta, and Mittal (1984) describe a branch and prune algorithm for solving  
39 UBQP which is designed to determine all local minimizing points, terminating with the global  
40 optimal solution revealed as the incumbent. Computational experience with random test  
41 problems with up to 125 variables is given.  
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46 Carter (1984) proposes a branch and bound algorithm for UBQP that first employs  
47 modified form of Cholesky factorization to transform an indefinite instance of UBQP into an  
48 equivalent positive definite form of the problem. Variable elimination based on hessian  
49 information is used accelerate search process. Computational experience on a variety of random  
50 test problems with various characteristics and with up to 30 variables is given.  
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56 Williams (1985) describes a branch and bound algorithm for UBQP that successfully  
57 solved a set of randomly generated test problems with up to 100 variables. The algorithm begins  
58 with a reduction procedure that obtains a good starting solution and subsequently uses the “roof  
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4 dual” to help guide the depth-first branch and bound search. Comparisons with other methods  
5 are given.  
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9 Barahona, Junger, and Reinelt (1989) describe an approach to solving UBQP that first  
10 reduces the problem to an equivalent instance of a max-cut problem. This, in turn, is solved by a  
11 linear programming-based branch and bound method. Constraints based on the cut polytope are  
12 used to improve node information and enhance the search process. Computational experience is  
13 reported on random problems up to size 100 variables along with comparisons with other  
14 methods.  
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20 Kalantari and Bagchi (1990) describe the adaptation of an algorithm for minimizing  
21 linearly constrained concave quadratic functions for the purpose of solving UBQP. Their  
22 method starts with a transformation to ensure the Q matrix is positive definite, giving an  
23 equivalent concave quadratic minimization problem. The authors then describe their branch-  
24 and-bound method where subproblems are defined by fixing a variable at zero or one and bounds  
25 are computed by minimizing a linear convex envelope over the feasible region of the  
26 subproblem. Computational testing is reported on random problems with up to 50 variables.  
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34 Pardalos and Rodgers (1990a) describe a branch-and-bound algorithm for solving UBQP  
35 that successfully solved a variety of test problems with up to 200 variables. The algorithm,  
36 which uses no multiplications or divisions, incorporates dynamic preprocessing techniques for  
37 fixing variables and heuristics for finding good starting points. In (1990b) the authors describe a  
38 parallel version of the algorithm implemented and tested on a hypercube architecture.  
39 Computational experience and an analysis of the speedup achieved are presented. Then in  
40 (1992), the authors describe a variation of their branch-and-bound algorithm designed for the  
41 UBQP representation of the maximum clique problem. Extensive computational experience is  
42 reported for alternative branching rules and data structures, in route to producing a specialized  
43 algorithm optimized for the maximum clique problem with up to 1000 vertices and 150,000  
44 edges.  
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55 Billionnet and Sutter (1994) describe a branch-and-bound algorithm for solving UBQP  
56 that successfully solved a large variety of random test problems with up to 100 variables. Their  
57 innovation was in the calculation of lower bounds to guide the search process. At each node in  
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4 the search tree, a lower bound is computed by combining information obtained from roof duality,  
5 a quadratic posiform associated with the directed cycles of an implications graph, and a  
6 component obtained from the induced posiform of degree 4. Computational experience is given  
7 comparing the method with other methods.  
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12 Palubeckis (1995) describes a branch-and-bound algorithm for solving UBQP utilizing  
13 heuristically generated subproblem solutions that are mapped onto the zero n-vector leading to  
14 transformed subproblems in the form of the original UBQP model. Special classes of polytope  
15 facets are employed in computing bounds used to guide the search process. Computational  
16 experience with random problems with up to 100 variables, and additional experience with  
17 several real problems having to do with printed circuit board design, illustrate the efficiency of  
18 the method and indicate that it compares favorably with other contemporary methods.  
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26 Helmsberg and Rendl (1998) describe a branch-and-bound algorithm for solving UBQP  
27 based on semidefinite relaxations and cutting planes to enhance the quality of the bounds  
28 produced. The semidefinite relaxations are solved by an interior point algorithm specialized for  
29 semi-definite programs. Computational experience is reported on a set of UBQP instances of the  
30 max-cut variety with up to 100 variables. While the approach was robust in that it was  
31 successful in solving the test problems attempted, run times were generally not competitive with  
32 other recently reported exact methods for UBQP.  
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40 Hansen, Jaumard, and Meyer (2000) describe an enhanced version of the branch-and-  
41 bound method of Pardalos and Rodgers (1990a) that led to favorable comparisons with the  
42 original algorithm on a standard set of problems with up to 100 variables. The new method  
43 employs improved bounds obtained by first transforming the problem to an equivalent posiform  
44 which yields tighter roof dual bounds as well as effective variable elimination test that efficiently  
45 guide the tree search process. The roof dual bounds are computed via a maximum flow  
46 algorithm.  
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53 Huang, Pardalos, and Prokopyev (2006) describe a depth-first branch-and-bound method  
54 that begins by first formulating an equivalent bi-level formulation of UBQP. This new  
55 formulation facilitates bounding procedures and pruning strategies, utilizing a gradient midpoint  
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4 method that proved to be effective in early testing. Computational experience with random test  
5 problems of various densities and with up to 60 variables is presented.  
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9         Pardalos, Prokopyev, and Busygin (2006) discuss the connections between discrete  
10 optimization and continuous optimization in general with a focus on formulations that embed the  
11 initial discrete domain into a larger continuous space. The authors then focus on the general  
12 UBQP model, indicating how a reformulation based on an appropriate diagonal perturbation,  
13 causing the Q matrix to be negative semidefinite, yields an equivalent continuous problem of  
14 minimizing a quadratic concave function over the unit hypercube. A discussion of this approach  
15 applied to the maximum clique problem is given.  
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23         Pan, Tan, and Jiang (2008) describe a continuous approach for solving UBQP based on  
24 the Fischer-Burmeister nonlinear complementarity function. Rather than employing relaxations  
25 and bounding information in a tree search scheme, the authors reformulate UBQP as a  
26 continuous problem with equilibrium constraints. In turn, the optimal solution to this model is  
27 found by a global continuation algorithm utilizing a strictly convex global smoothing function  
28 and solving a sequence of unconstrained minimization problems. Computational experience is  
29 reported with random problems with up to 1000 variables, indicating the effectiveness of this  
30 approach.  
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38         Gueye and Michelon (2009) present a general framework for constructing linearizations  
39 of UBQP which, in turn, can be solved in principle by standard optimizers for mixed integer  
40 linear programs. The framework, which contains existing linearization methods in the literature  
41 as special cases, consists of decomposing the objective function into component matrices,  
42 identifying a complete linear representation of the polytope for each component, and then adding  
43 constraints that link the components together. A new linearization, derived from the general  
44 framework, is described and computational comparisons are given with existing methods  
45 illustrating the potential of the new approach.  
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53         Dinh, Canh, and Thi (2010) discuss a new continuous approach for solving UBQP that is  
54 based on DC (difference of convex functions) programming. In their approach, principles of DC  
55 programming are used to develop a local optimization algorithm (DCA) that solves a finite  
56 number of linear programs leading to a locally optimal solution. Globally optimal solutions are  
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4 produced by embedding DCA in a branch-and-bound algorithm. Computational experience on  
5 problems from the literature with up to 100 variables, along with comparison with other  
6 methods, is given.  
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10 Mauri and Lorena (2011) present a new algorithm for solving UBQP based on  
11 Lagrangean decompositions. Their method starts with a linearization of UBQP represented by a  
12 graph. In turn, the graph is partitioned into clusters of vertices forming a dual problem that is  
13 solved by a subgradient algorithm. Clusters are formed using the well-known METIS heuristic  
14 and the linear Lagrangean subproblems are solved using CPLEX. In (2012a) the authors present  
15 a column generation alternative to the subgradient algorithm leading to performance  
16 improvement. In (2012b) the authors present and test further enhancements to their column  
17 generations approach for solving UBQP. Throughout all, computational experience on standard  
18 UBQP test problems with up to 500 variables is given. Comparisons with other decomposition-  
19 based methods are given indicating the potential of the procedures proposed here.  
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30 Li, Sun, and Liu (2012) present a new algorithm for solving UBQP based on the inherent  
31 geometric properties of the minimum circumscribed sphere containing the ellipsoidal contour of  
32 the objective function. Based on these properties, effective bounding information as well as new  
33 procedures for optimally fixing variables are derived. In addition, this geometric approach led to  
34 new optimality conditions for UBQP. The new bounding techniques and variable fixing  
35 conditions were combined in a branch-and-bound method and tested on standard problems with  
36 up to 200 variables. Comparisons drawn with other recent methods in the literature indicate the  
37 attractiveness of the method proposed.  
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45 In addition to the exact methods surveyed above from the literature, we point out that  
46 several commercial methods, based on branch-and-cut techniques, are now available and hold  
47 considerable promise for directly optimizing moderate sized instances of UBQP. See for  
48 example the paper by Billionnet & Elloumi (2007) which reports on the use of the branch-and  
49 cut quadratic integer optimizer available from CPLEX.  
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### 55 3.2 Heuristic and Metaheuristic Methods 56

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58 The NP-hard nature of UBQP along with its application potential has motivated a large  
59 number of papers in recent years describing various heuristic methods for quickly finding high  
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4 quality solutions to medium to large sized problem instances. Although a few of these methods  
5 are simple enough to qualify as heuristics, those that generate the best solutions are metaheuristic  
6 procedures that incorporate compound strategies considerably more advanced than in the basic  
7 heuristics. These methods are surveyed below:  
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12 Boros, Hammer and Sun (1989) develop a Devour Digest Tidy-up (DDT) procedure to  
13 rapidly obtain a solution to UBQP. Based on the posiform expression of UBQP's objective  
14 function, the proposed method includes devour, digest and tidy-up phases. The devour phase  
15 identifies a term with the largest coefficient and sets it to 0 in terms of minimization. The digest  
16 phase draws logical conclusions for the items from the devour phase. The tidy-up phase finally  
17 substitutes the logical consequences previously derived into the current quadratic function.  
18 Computational experience indicates the effectiveness of the method proposed, in particular on  
19 problems of low density.  
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28 Glover, Kochenberger, and Alidaee (1998) propose an adaptive memory tabu search  
29 algorithm, which incorporates a strategic oscillation scheme to enable the search to go beyond  
30 the local optimum obtained by constructive and destructive phases. A key feature of this method  
31 lies in the use of a *critical event memory*, that collects recency and frequency information from  
32 critical events (moves that causes the solution values to decrease), to guide the oscillation  
33 process. Another feature lies in the use of adaptive oscillation depths. Extensive computational  
34 experience discloses that the proposed method outperforms the best exact and heuristic methods  
35 previously reported in the literature in terms of speed and solution quality.  
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44 Glover, Kochenberger, Alidaee and Amini (1999) describe an enhanced version of their  
45 previous adaptive memory tabu search algorithm. A simple but effective scheme is proposed for  
46 accelerating the evaluation of moves and for updating associated problem information. In  
47 addition, methods for generating high quality initial solutions and for creating additional trial  
48 solutions at critical events are also introduced. Computational experience with up to 1000  
49 variables reveals this enhanced version can produce high quality solutions within several  
50 minutes.  
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57 Beasley (1998) adapts tabu search and simulated annealing to solving UBQP. The tabu  
58 search implementation incorporates a strategy in which, once an improved solution is found, a  
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4 simple local search is successively employed to perform moves irrespective of their tabu status.  
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6 Contrary to the tabu search procedure, a local search is applied only at the end of the simulated  
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8 annealing process. Computational comparisons indicate that their tabu search generally performs  
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10 better than simulated annealing for small and medium instances but worse for large instances.  
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13 Alkhamis, Hasan, and Ahmed (1998) present a simulated annealing based heuristic with a  
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15 well-selected cooling schedule. Tested on several hundred test problems, the proposed heuristic  
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17 outperforms several algorithms based on bounding techniques, in particular with respect to  
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19 computational time. Additional analysis shows that initial solutions and the matrix density have  
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21 limited influence on the effectiveness of the simulated annealing algorithm.  
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24 Merz and Freisleben (1999) devise a hybrid genetic algorithm, in which a simple local  
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26 search is incorporated into the traditional genetic algorithm. The crossover operator is a variant  
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28 of uniform crossover, requiring the generated offspring solutions to have the same hamming  
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30 distance from the parents. The population updating criterion refers to the quality of solutions,  
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32 assuring that each solution occurs only once in the population, as customarily done in scatter  
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34 search methods. A diversification component is launched when the average hamming distance of  
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36 the population drops below a threshold or the population is not updated for more than 30  
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38 consecutive generations. Computational experience shows that a simple genetic algorithm is  
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40 sufficient to find best known results for problem instances with less than 200 variables but for  
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42 those with a large number of variables, local search is needed for attaining high quality solutions.  
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45 Amini, Alidaee, and Kochenberger (1999) present a scatter search approach, which  
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47 consists of a diversification generation method, a solution improvement method, a reference set  
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49 update method, a subset generation method and a solution combination method. The  
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51 diversification generation method systematically generates a collection of diverse trial solutions  
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53 based on a seed solution by setting an incremental parameter that determines which bits of the  
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55 seed solution should be flipped. The improvement method performs a compound move that  
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57 sequentially cycles among three types of candidate moves until no attractive move can be  
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59 identified. The reference set update method replaces solutions in the reference set with new  
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61 candidate solutions using the quality measurement. The solution combination method uses linear  
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63 combination of solutions in a subset derived from the subset generation method to produce new  
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65 solutions. Since some variables may receive fractional values in the solution obtained in the

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4 linear combination, a rounding procedure is employed to recover integer values. Experiments on  
5 three classes of problems show the attractiveness of the proposed method.  
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9 Lodi, Allemand, and Liebling (1999) present an evolutionary method for solving UBQP.  
10 The proposed algorithm is characterized by the following features. First, a preprocessing phase is  
11 applied to fix certain variables at their optimal values and keep them unchanged during each  
12 successive round of local search, hence resulting in a reduced problem scale. Second, a local  
13 search procedure that alternates between construction phase and destructive phases is used to get  
14 an improved solution. Finally, a uniform crossover operator is used to produce offspring  
15 solutions, where variables with common values in parental solutions are temporarily fixed in this  
16 round of local search. Computational experience on problem instances with up to 500 variables  
17 is given. A further analysis demonstrates that the preprocessing phase is effective for small  
18 problem instances but is unable to appreciably reduce the problem size for large ones.  
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28 Katayama, Tani, and Narihisa (2000) propose a genetic local search algorithm for solving  
29 UBQP. Their local search procedure integrates *1-flip* moves dedicated to going into new good  
30 search area and *k-flip* moves dedicated to solution improvement. A traditional uniform crossover  
31 and a simple mutation operator are joined to generate a suitable offspring solution. A  
32 diversification/restart strategy is incorporated to maintain a diversified population. Tests on large  
33 problem instances indicate the effectiveness of the proposed algorithm.  
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40 Katayama and Narihisa (2001) present a simulated annealing algorithm with an  
41 innovative use of multiple annealing processes to enhance the search. Each annealing process  
42 takes the best solution found in the previous annealing process as the initial solution and employs  
43 a different initial temperature. Experimental results demonstrate the performance of the proposed  
44 algorithm, especially for large instances with 2500 variables.  
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50 Merz and Freisleben (2002) describe a greedy heuristic and two local search algorithms  
51 based upon *1-flip* and *k-flip* neighborhoods. The greedy construction procedure starts from a  
52 solution with all variables assigned to 0.5 (the so called third state) and each constructive step  
53 picks a variable with probability proportional to the gain value when changing the variable's  
54 value from 0.5 to 0 or 1. Each iteration of the *1-flip* local search proceeds to the neighbor  
55 solution with the best solution quality. The *k-flip* local search borrows the idea from the Lin-  
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4 Kernighan algorithm of Kernighan and Lin (1972) for solving the graph partitioning problem to  
5 efficiently reduce the neighborhood exploration. Each *k-flip* move consists in repeating  
6 performing the best *1-flip* move until all *1-flip* moves are performed and picking the best from  
7 the resulting solutions. Computational comparisons disclose the superiority of the multistart *k-*  
8 *flip* local search combined with randomized greedy initial solutions.  
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14 Glover, Alidaee, Rego, and Kochenberger (2002) propose several one-pass heuristics to  
15 advance the Devour Digest Tidy-up (DDT) method of Boros, Hammer and Sun (1989). Based on  
16 the hypothesis that setting multiple variables with value 1 or 0 in a pass may lead to worse  
17 performance, the idea is to guarantee only one variable gets the implied assignment in each pass.  
18 The proposed one-pass heuristics differ in strategies for evaluating contributions of variables.  
19 Computational experience indicates that the method outperforms the DDT method but no single  
20 one-pass heuristic dominates the others on every problem instance.  
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28 Palubeckis and Tomkevicius (2002) present a greedy random adaptive search procedure  
29 (GRASP) which cycles between a construction phase and a local search phase. Each step in the  
30 construction phase picks a variable from a candidate list with probability proportional to the gain  
31 value of this variable, where the candidate list is composed of a certain number of variables with  
32 the largest gain values, calculated according to a specific gain function. The local search phase  
33 implements a simple ascent algorithm. Two enhanced versions are tested, which result by  
34 replacing local search with tabu search and by combining a classic random restarting procedure  
35 with tabu search. Computational comparisons illustrate the merit of incorporating greedy  
36 construction based initial solutions and tabu search.  
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45 Palubeckis (2004) examines five multistart tabu search strategies dedicated to the  
46 construction of an initial solution. The first multistart strategy produces a new initial solution in  
47 a random way. The second identifies a candidate set of variables whose values are prone to  
48 change when moving from the current solution to an optimal one and then applies a steepest  
49 ascent algorithm where variables not included in this candidate set are fixed at specific values.  
50 The third multistart strategy is the same as the constructive phase proposed in Palubeckis, 2002.  
51 The fourth uses a set of elite solutions to calculate the probability of each variable being assigned  
52 value 1. If the probability for a given variable is larger than 0.5, then this variable is assigned to  
53 be 1 in the constructed solution; otherwise it is assigned to be 0. The last multistart strategy uses  
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4 a perturbation scheme of changing the problem instance at hand, followed by a short run of tabu  
5 search on the modified instance. Extensive comparisons on problem instances with up to 7000  
6 variables demonstrate the algorithm using the second multistart strategy performs better than the  
7 other proposed alternatives.  
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12 Merz and Katayama (2004) conduct landscape analysis and observe that (1) local optima  
13 of the UBQP problem instances are concentrated in a small fraction of the search space; (2) the  
14 fitness of local optima and the distance between local optima and the global optimum are  
15 correlated. Based on the observations, they propose a memetic algorithm in which an innovative  
16 variation operator is used to generate an offspring solution and the *k-flip* local search proposed in  
17 Katayama, Tani, and Narihisa (2000) is used to improve solution quality. The variation operator  
18 introduces new alleles not contained in both parents by referring to the move gain of performing  
19 *l-flip* moves, avoiding the rediscovery of local optima already extensively visited. Comparisons  
20 with other algorithms demonstrate the effectiveness of the proposed algorithm.  
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30 Boros, Hammer, and Tavares (2006) present several preprocessing techniques to simplify  
31 the UBQP problem. The purpose of the preprocessing simplification is to provide several  
32 features, including lower bounds for the minimum of the objective function, optimal assignments  
33 for some variables, and binary relations between the values of certain pairs of variables and  
34 subproblems decomposed from the original problem. The simplification is achieved by using  
35 basic techniques such as first-order derivatives, second-order derivatives or roof-duality, and by  
36 using integrative techniques that combine the conclusions derived from the basic techniques.  
37 Computational experience on numerous problem classes shows the value of the proposed  
38 preprocessing techniques.  
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47 Palubeckis (2006) presents an iterated tabu search algorithm which uses a dedicated  
48 perturbation mechanism to enhance the high-quality solution obtained by the tabu search  
49 procedure. Each step of the perturbation constructs a candidate list of a limited size consisting of  
50 variables with largest *l-flip* move gains with regard to this current solution, from which a  
51 variable is randomly selected and flipped to complement the value of this variable. The current  
52 solution is thus updated and the next perturbation step continues until the number of perturbed  
53 variables reaches the specified number. Comparisons with state-of-the-art algorithms disclose the  
54 competitiveness of the proposed algorithm in spite of its simplicity.  
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4 Boros, Hammer, and Tavares (2007) present a local search scheme for solving UBQP.  
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6 Starting from an initial solution, each iterative step constructs a candidate set from which a  
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8 variable is picked and its value is changed to its complement, thus moving to the next solution.  
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10 This iterative procedure repeats until the candidate set becomes empty. Based on the above  
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12 scheme, they investigate five initialization methods, two candidate set construction methods and  
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14 four variable selection methods, thus reaching up to 40 local search alternatives. Experiments on  
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16 multiple benchmark instances indicate that the local search alternative combining the following  
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18 methods achieves the best performance. The initial method assigns each variable with a  
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20 fractional value equaling to the proportion of the sum of all the positive entries of the matrix in  
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22 the sum of the absolute value of each entry of the matrix. The candidate set construction method  
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24 constructs a candidate set consisting of variables that yield an improvement in the current  
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26 solution by flipping its value regardless of whether or not it was already flipped in the previous  
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28 iteration. The variable selection method selects from the candidate set the variable with the  
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30 largest improvement to the current solution.

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32 Glover, Lu, and Hao (2010) present a diversification-driven tabu search algorithm, which  
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34 alternates between a basic tabu search procedure and a memory-based perturbation strategy  
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36 guided by a long-term memory. Three memory structures are introduced in the perturbation  
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38 strategy: (1) a flipping frequency vector to record the number of times a variable has been  
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40 flipped from the initial iteration until the current iteration; (2) an elite set of solutions to record a  
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42 certain number of best local optimal solutions; (3) a consistency vector to count the times each  
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44 variable is assigned a given value in the set of elite solutions. Based on the memory information,  
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46 the perturbation operator modifies an elite solution by favoring variables with low flipping  
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48 frequency and high consistency to flip. Comparisons drawn with several algorithms proposed by  
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50 Palubeckis (2004, 2006) disclose the superiority of this algorithm.

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52 Lu, Glover, and Hao (2010) present a hybrid metaheuristic approach which has the  
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54 following features. First, it combines a traditional uniform crossover operator with a  
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56 diversification guided path relinking operator to guarantee the quality and diversity of an  
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58 offspring solution. Second, it defines a new distance by reference to variable's importance and  
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60 employs a quality-and-distance criterion to update the population. Finally, a tabu search  
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62 procedure is responsible for intensified examination around the offspring solutions.  
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4 Computational comparisons with best performing algorithms indicate the effectiveness of this  
5 hybrid algorithm.  
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9 Lu, Hao, and Glover (2010) develop a hybrid genetic tabu search with multi-parent  
10 crossover to solve UBQP. The proposed algorithm jointly uses traditional uniform crossover and  
11 logic multi-parent combination operators to generate diversified offspring solutions.  
12 Computational experience is given showing the competitiveness of the proposed algorithm.  
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17 Cai, Wang, Yin, and Zhou (2011) present a memetic clonal selection algorithm with  
18 estimation of distribution algorithm (EDA) guided vaccination for solving UBQP. The proposed  
19 algorithm adopts EDA vaccination, fitness uniform selection scheme and adaptive tabu search to  
20 overcome the deficiencies of traditional clonal selection algorithm. Experimental comparisons  
21 indicate the tabu search algorithm enhances the performance of the clonal selection algorithm.  
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27 Shylo and Shylo (2011) develop a global equilibrium search which performs multiple  
28 temperature cycles. Each temperature cycle includes an initial solution generation phase and a  
29 tabu search phase. The method to generate an initial solution employs historical information to  
30 determine the probability that a variable receives the value 1. The tabu search procedure requires  
31 that each admissible move leads to a solution with hamming distance to a reference set  
32 surpassing a distance threshold. Computational comparisons with several algorithms indicate the  
33 attractiveness of the proposed algorithm.  
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41 Hanafi, Rebai, and Vasquez (2013) devise five alternative DDT heuristics based on  
42 different representations of the BQO formulation. DDT1 to DDT4 respectively have standard,  
43 posiform, bi-form and negaform representations and DDT5 has a posiform representation  
44 combined with a one-pass mechanism. One obvious difference between the DDT alternatives  
45 proposed here and those proposed by Boros et al. (1989) and Glover et al. (2002) lies in the use  
46 of an *r-flip* local search procedure to improve solutions obtained by DDT constructions.  
47 Extensive tests on small, medium and large benchmark instances disclose that (1) DDT3 with the  
48 bi-form representation generally produces the best results for medium and large instances; (2) the  
49 *r-flip* local search contributes to significant improvement of the results of the proposed DDT  
50 methods with only a slight increase of time consumption.  
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4 Wang, Zhou, and Yin (2011) present a tabu Hopfield neural network with an estimation  
5 of distribution algorithm (EDA). The cooperation between long term memory of EDA with the  
6 short term memory of tabu search prevents the network from becoming trapped in local optima.  
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8 Computational testing indicates the superiority of the proposed algorithm compared to other  
9 Hopfield neural network based algorithms.  
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14 Lu, Glover, and Hao (2011) study neighborhood union and token-ring search methods  
15 within a tabu search algorithm. They focus on two neighborhoods, N1 consisting of *1-flip* moves  
16 and N2 consisting of a chosen subset of *2-flip* moves. The neighborhood union includes the  
17 strong neighborhood union that picks each move from both N1 and N2 and the selective  
18 neighborhood union that picks a move from N1 with probability  $p$  and N2 with probability  $1-p$ .  
19 The token ring search continuously performs move in N1 until no improvement is possible and  
20 then switches to perform move in N2 to continue the search. Computational comparisons reveal  
21 the superiority of the token ring search over the neighborhood union.  
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30 Wang, Lu, Glover and Hao (2012) present two path relinking algorithms, which are  
31 composed of a reference set construction method, a tabu search based improvement method, a  
32 reference set update method, a relinking method and a path solution selection method. The  
33 proposed algorithms differ from each other mainly on the way they generate the path, one  
34 employing a greedy strategy and the other employing a random strategy. Extensive  
35 computational experience and comparisons with several state-of-the-art algorithms highlight the  
36 attractiveness of the proposed algorithms in terms of both solution quality and computational  
37 efficiency.  
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45 Wang, Lu, Glover and Hao (2012) propose a simple GRASP-Tabu Search algorithm  
46 working with a single solution and an enhanced version by combining GRASP-Tabu Search  
47 algorithm with a population management strategy based on an elite reference set. In the basic  
48 version, the initial solution is constructed according to a greedy random construction heuristic. In  
49 the enhanced version, a new solution is reconstructed by first inheriting parts of the good  
50 assignments of one elite solution to form a partial solution and then completing the remaining  
51 parts as the basic version does. Experimental tests on a large range of both random and  
52 structured problem instances disclose that the proposed algorithms, in particular the enhanced  
53 version, yield very competitive outcomes.  
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4 Wang, Lu, Glover, and Hao (2012) present a backbone guided tabu search algorithm  
5 which alternates between a basic tabu search procedure and a variable fixing/freeing phase based  
6 on identifying strongly determined variables. While the tabu search phase ensures the  
7 exploitation of a search space, the variable fixing (freeing) phase dynamically enlarges (reduces)  
8 the backbone of assigned values that launches the tabu search exploration. Experiments show  
9 that the proposed algorithm obtains highly competitive outcomes in comparison with the  
10 previous best known results from the literature. A direct comparison with the underlying tabu  
11 search procedure confirms the merit of incorporating backbone information.  
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20 As indicated in the papers of this section, our ability to efficiently solve large instances of  
21 UBQP by heuristic means has grown substantially in recent years. It is common now for authors  
22 to report computational experience on problems with 7,000 – 10,000 variables. Note that the set  
23 partitioning application discussed in section 2.3 by Lewis, et. al. (2008) reported computational  
24 experience on problems up to 15,000 variables.  
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#### 29 30 **4.0 Key Theoretical Results**

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32 By far the majority of the papers in the literature related to UBQP are primarily devoted  
33 to applications or various solution schemes, either exact or heuristic in nature. As a result, our  
34 priority in this paper has been to focus our survey on applications and solution methodologies.  
35 Many of the articles surveyed in section 3 above, however, contain a discussion of the theoretical  
36 results relevant to the method being put forth. That is, these papers are *mainly* about the method  
37 at hand but may also contain a discussion of underlying theory. As a result, we've not explicitly  
38 highlighted theoretical issues but rather left them to be discovered, as might be appropriate, as  
39 part of the articles on applications and solutions methods surveyed. Nonetheless, there are a few  
40 recent papers in the literature of particular note focused on theoretical issues pertaining to  
41 UBQP. It is these papers that we highlight here in this section.  
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50 Carraresi, Farinaccio and Malucello (1999) present an exact algorithm for testing the  
51 optimality of a given solution for a quadratic 0-1 unconstrained problem. Their method, based on  
52 necessary and sufficient conditions introduced by Hirriart-Urruty for general convex problems,  
53 expands their earlier work (1995) which was an approximation scheme for testing solutions.  
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57 Beck and Teboulle (2000) characterize global optimal solutions for UBQP as well as  
58 discussing the relationship between optimal solutions to UBQP and the optimal solutions of its  
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4 continuous relaxation. They derive a sufficient optimality condition which guarantees that a  
5 given feasible point is a global optimal for UBQP as well as a necessary global optimality  
6 condition.  
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10 Jeyakumar, Rubinov, and Wu (2007) examine the relationship between the global  
11 optimality of nonconvex constrained optimization and Lagrange multiplier conditions,  
12 establishing sufficient as well as necessary conditions for global optimality for general quadratic  
13 minimization problems with quadratic constraints. This analysis led, as a special case, to new  
14 sufficient and necessary global optimality conditions for UBQP that are sharper than those given  
15 earlier by Beck and Teboulle.  
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19 Xia (2009), by analyzing local sufficient optimality conditions, also extended the Beck  
20 and Teboulle results by developing tighter sufficient optimality conditions. In addition, without  
21 making the positive-semidefinite assumption, Xia examines the relationship between local/global  
22 minimizers of UBQP and the KKT points of the continuous relaxation, further extending  
23 previous results in the literature.  
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27 Gao and Ruan (2010) present a discussion of canonical duality theory, designed in  
28 general for a wide class of nonconvex/nonsmooth/discrete problems. The authors show how this  
29 duality theory can be adapted for the quadratic case with binary constraints. Conditions are  
30 given that allow instances of UBQP to be converted into smooth concave maximization dual  
31 problems over a closed convex feasible region without a duality gap. Finally, the relationship  
32 between canonical duality theory and semi-definite programming for UBQP is discussed.  
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36 Zheng, Sun, Li, and Xu (2012) present new sufficient conditions for verifying zero  
37 duality gap in nonconvex constrained quadratic programs and then show how the results  
38 specialize for UBQP. In related work, Sun, Liu, and Gao (2012) investigate the duality gap  
39 between UBQP and its semi-definite programming relaxation. Making the connection between  
40 the duality gap and the cell enumerations of hyperplane arrangement in discrete geometry,  
41 estimates of the duality gap can be derived, yielding improved lower bounds for UBQP.  
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45 We note that there are several theoretical papers in the literature on the constrained  
46 version of UBQP that don't explicitly consider the pure UBQP model but are nonetheless  
47 relevant to our work in that UBQP is a special case of the constrained cases considered.  
48 Notably, Pinar (2004) gives a discussion of sufficient global optimality conditions for the  
49 problem of minimizing a quadratic function in binary variables subject to equality quadratic  
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4 constraints. Lu, et. al., (2011) presents a discussion of KKT conditions and conic relaxations to  
5 develop sufficient conditions that generalize known positive semi-definiteness results for finding  
6 globally optimal solutions for the problem of minimizing a UBQP subject to inequality quadratic  
7 constraints. Finally, Li (2012) presents an extension of Pinar's global optimality conditions for  
8 the quadratic equality constrained case along with presenting conditions enabling global  
9 optimality to be assessed by checking the positive semi-definiteness of a related matrix.  
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## 16 **5.0 Summary & Conclusions**

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19 Interest in UBQP has grown substantially in recent years as researchers have discovered  
20 the remarkable ability of this simple model form to represent a wide variety of combinatorial  
21 problems along with its computational challenge, particularly as model sizes have increased.  
22 Due to its NP-hard nature, methods capable of producing exact solutions are limited to modest  
23 sized applications, giving way to modern heuristic methods for larger models. Even today, exact  
24 methods appear to be limited to a few hundred variables. In an effort to realize the application  
25 potential of UBQP as model size scales to higher levels, most research is focused on  
26 metaheuristic methods of one kind or another. The results are encouraging: Articles in the 80s  
27 were reporting on solving problems with 100-200 variables while more recent articles are  
28 reporting on problems with up to 15,000 variables. To a large extent, this growth in performance  
29 is due to advances in both algorithm design and computer hardware.  
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41 Successfully moving to the next order of magnitude in terms of model size will require  
42 creative schemes for handling very large Q matrices along with improved algorithmic methods.  
43 Various partitioning and multi-level methods hold particular promise here but the door is open  
44 for other innovations as well. Advances in computer performance, both in terms of storage and  
45 speed, can also be expected to lend a hand in allowing larger applications. Moreover,  
46 developments in the area of quantum computing, as illustrated by the work by Neven, Rose, and  
47 Macready (2008), represent emerging technologies with a potential for solving combinatorial  
48 problems as represented by UBQP. Future papers will reveal which of these research areas, or  
49 indeed, if some other approach, will contribute to facilitating solutions to UBQP as application  
50 size continues to scale upward.  
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